Software Architecture an informal introduction

David Schmidt

Kansas State University

www.cis.ksu.edu/~schmidt

Outline

- **1. Components and connectors**
- 2. Software architectures
- 3. Architectural analysis and views
- 4. Architectural description languages
- 5. Domain-specific design
- 6. Product lines
- 7. Middleware
- 8. Model-driven architecture
- 9. Aspect-oriented programming
- **10. Closing remarks**

1. Components and connectors

Motivation for software architecture

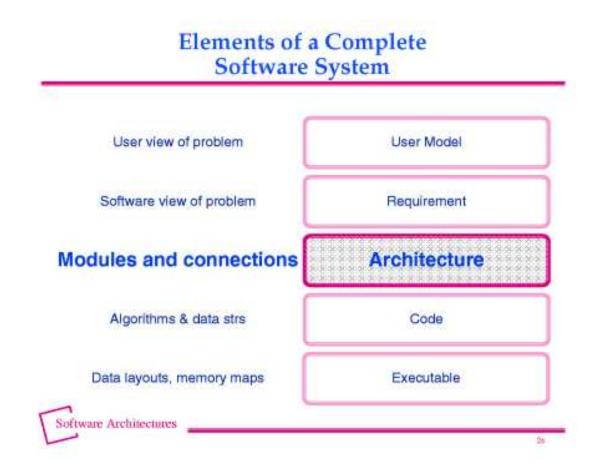
We use already architectural idioms for describing the structure of complex software systems:

- "Camelot is based on the *client-server model* and uses remote procedure calls both locally and remotely to provide communication among applications and servers." [Spector87]
- "The easiest way to make the canonical sequential compiler into a concurrent compiler is to *pipeline* the execution of the compiler phases over a number of processors." [Seshadri88]
- "The ARC network follows the general network architecture specified by the ISO in the Open Systems Interconnection Reference Model." [Paulk85]

Reference: David Garlan, Architectures for Software Systems, CMU, Spring 1998. http://www.cs.cmu.edu/afs/cs/project/tinker-arch/www/html/index.html

Architectural description has a natural position in system design and implementation

A slide from one of David Garlan's lectures:



Reference: David Garlan, *Architectures for Software Systems*, CMU, Spring 1998. http://www.cs.cmu.edu/afs/cs/project/tinker-arch/www/html/index.html

Hardware architecture

There are standardized descriptions of computer hardware architectures:

- *RISC* (reduced instruction set computer)
- pipelined architectures
- *multi-processor architectures*

These descriptions are well understood and successful because

(i) there are a relatively small number of design components

(ii) large-scale design is achieved by replication of design elements

In contrast, software systems use a huge number of design components and scale upwards, not by replication of existing structure, but by adding more distinct design components.

Reference: D. E. Perry and A. L. Wolf. Foundations for the Study of Software Architectures. *ACM SIGSOFT Software Engineering Notes*, October 1992.

Network architecture

Again, there are standardized descriptions:

- star networks
- ring networks
- manhattan street (grid) networks

The architectures are described in terms of *nodes* and *connections*. There are only a few standard topologies.

In contrast, software systems use a wide variety of topologies.

Classical architecture

The architecture of a building is described by

- *multiple views*: exterior, floor plans, plumbing/wiring, ...
- ♦ architectural styles: romanesque, gothic, ...
- style and engineering: how the choice of style influences the physical design of the building
- style and materials: how the choice of style influences the materials used to construct (implement) the building.

These concepts also appear in software systems: there are

- (i) views: control-flow, data-flow, modular structure, behavioral requirements, ...
- (ii) styles: pipe-and-filter, object-oriented, procedural, ...
- (iii) *engineering*: modules, filters, messages, events, ...
- (iv) *materials*: control structures, data structures, ...

A crucial motivating concept: connectors

The emergence of networks, client-server systems, and OO-based GUI applications led to the conclusion that

components can be connected in various ways

Mary Shaw stressed this point:

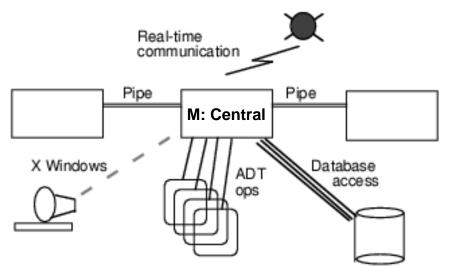


Figure 2: Revised architecture diagram with discrimination among connections

Reference: Mary Shaw, Procedure Calls are the Assembly Language of Software Interconnections: Connectors Deserve First-Class Status. Workshop on Studies of Software Design, 1993.

Shaw's observations

Connectors are forgotten because (it appears that) there are no codes for them.

But this is because the connectors must be coded in the same language as the components, which confuses the two forms.

Different forms of low-level connection (synchronous, asynchronous, peer-to-peer, event broadcast) are fundamentally different yet are all represented as procedure (system) calls in programming language.

Connectors can (and should?) be coded in languages different from the languages in which components are coded (e.g., unix pipes). **Components** — compilation units (module, data structure, filter) — are specified by *interfaces*.

Connectors — "hookers-up" (RPC (Remote Procedure Call), event, pipe) — mediate communications between components and are specified by *protocols*.

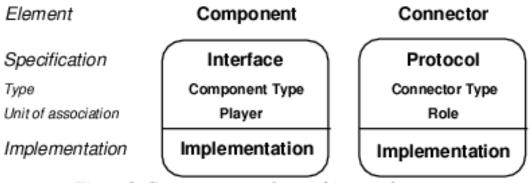


Figure 3: Gross structure of an architecture language

Example:

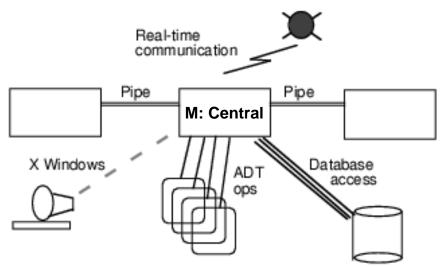
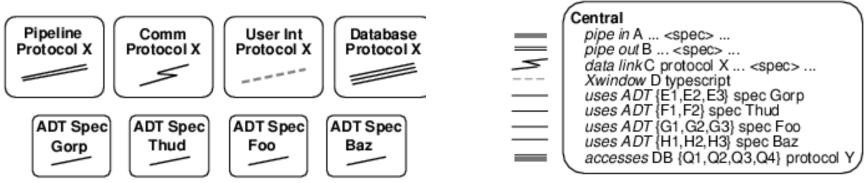


Figure 2: Revised architecture diagram with discrimination among connections



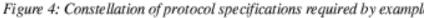


Figure 5: Interface specification of central component, referring to protocols

Interface **Central** is different from a Java-interface; it lists the "players" — inA, outB, linkC, Gorp, Thud, ... (connection points/ ports/ method invocations) — that use connectors.

The connector's *protocol* lists

(i) the types of component interfaces it can "mediate";

(ii) orderings and invariants of component interactions;

(iii) performance guarantees.

Example: Shaw's description of a unix pipe:

Pipe Connector

Informal Description: The Unix abstraction for pipe, i.e. a bounded queue of bytes that are produced at a source and consumed at a sink. Also supports interactions between pipes and files, choosing the correct Unix implementation.

Icon: pipe section

Properties: PipeType, the kind of Unix pipe. Possible values Named, Unnamed

Roles: Source

Description: the source end of the pipe

Accepts player types: StreamOut of component Filter; ReadNext of component SeqFile Properties: MinConns, minimum number of connections. Integer values, default 1 MaxConns, maximum number of connections. Integer values, default 1

Sink

Description: the sink end of the pipe Accepts player types: StreamIn of component Filter; WriteNext of component SeqFile Properties: MinConns, MaxConns, as for Source

Reference: M. Shaw, R. DeLine, and G. Zelesnik. Abstractions and Implementations for Architectural Connections. In 3d. Int. Conf. on Configurable Distributed Systems, Annapolis, Maryland, May 1996.

Connectors can act as

- communicators: transfer data between components (e.g., message passing, buffering)
- *mediators:* manage shared resource access between components (e.g., reader/writer policies, monitors, critical regions)
- coordinators: define control flow between components (e.g., synchronization (protocols) between clients and servers, event broadcast and delivery)
- adaptors: connect mismatched components (e.g., a pipe connects to a file rather than to a filter)

Perhaps you have written code for a bounded buffer or a monitor or a protocol or a shared, global variable — you have written a connector!

Connectors can facilitate

- reuse: components from one application are inserted into another, and they need not know about context in which they are connected
- evolution: components can be dynamically added and removed from connectors
- heterogenity: components that use different forms of communication can be connected together in the same system

A connector should have the ability to handle limited *mismatches* between connected components, via information reformatting, object-wrappers, and object-adaptors, such that the component is not rewritten — the connector does the reformatting, wrapping, adapting.

If connectors are crucial to systems building, why did we take so long to "discover" them? One answer is that components are "pre-packaged" to use certain connectors:

Component t	YPE COMMON TYPES OF INTERACTION
Module	Procedure call, data sharing
Object	Method invocation (dynamically bound procedure call)
Filter	Data flow
Process	Message passing, remote procedure call
	various communication protocols, synchronization
Data file	Read, write
Database	Schema, query language
Document	Shared representation assumptions

But "smart" connectors make components simpler, because the coding for interaction rests in the connectors — not the components.

The philosophy, **system = components + connectors** was a strong motivation for a theory of software architecture.

Reference: M. Shaw and D. Garlan. Formulations and Formalisms in Software Architecture. *Computer Science Today: Recent Trends and Developments* Jan van Leeuwen, ed., Springer-Verlag LNCS, 1996, pp. 307-323.

2. Software Architecture

What is a software architecture? (Perry and Wolf)

A software architecture consists of

- elements: processing elements ("functions"), connectors ("glue" procedure calls, messages, events, shared storage cells), data elements (what "flows" between the processing elements)
- 2. *form*: properties (constraints on elements and system) and relationship (configuration, topology)
- 3. *rationale*: philosophy and pragmatics of the system: requirements, economics, reliability, performance

There can be "views" of the architecture from the perspective of the process elements, the data, or the connectors. The views might show static and dynamic structure.

Reference: D. E. Perry and A. L. Wolf. Foundations for the Study of Software Architectures. *ACM SIGSOFT Software Engineering Notes*, October 1992.

Architectural Styles (patterns)

- 1. *Data-flow systems:* batch sequential, pipes and filters
- 2. *Call-and-return systems:* main program and subroutines, hierarchical layers, object-oriented systems
- 3. *Virtual machines:* interpreters, rule-based systems
- 4. *Independent components:* communicating systems, event systems, distributed systems
- 5. *Repositories (data-centered systems):* databases, blackboards
- 6. and there are many others, including *hybrid* architectures

The *italicized* terms are the styles (e.g., *independent components*); the roman terms are architectures (e.g., communicating system). There are specific instances of the architectures (e.g., a client-server architecture is a distributed system). But these notions are not firm!

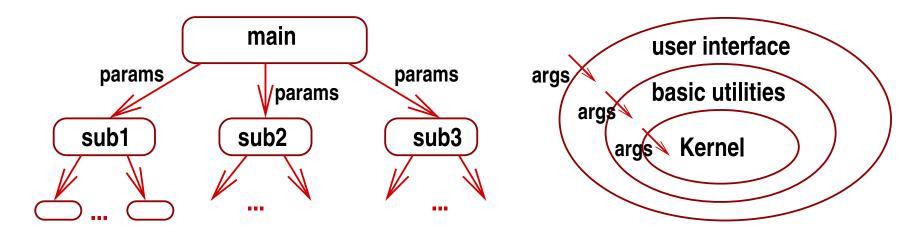
Data-flow systems: Batch-sequential and Pipe-and-filter

text Scan	tokens	tree GenCode code
	Batch sequential	Pipe and filter
Components:	whole program	filter (function)
Connectors:	conventional input-output	pipe (data flow)
Constraints:	components execute to completion, consuming entire input, producing entire output	data arrives in incre- ments to filters

Examples: Unix shells, signal processing, multi-pass compilers

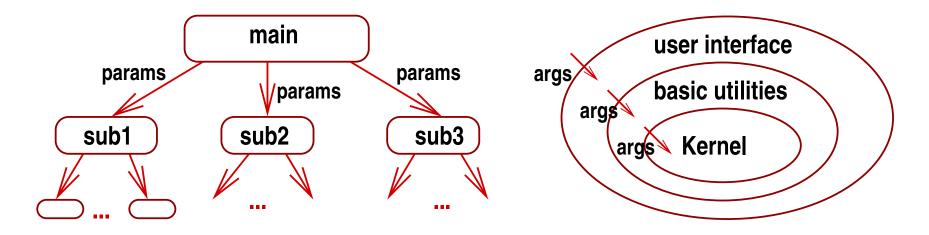
Advantages: easy to unplug and replace filters; interactions between components easy to analyze. *Disadvantages:* interactivity with end-user severely limited; performs as quickly as slowest component.

Call-and-return systems: subroutine and layered



	Subroutine	Layered
Components:	subroutines ("servers")	functions ("servers")
Connectors:	parameter passing	protocols
Constraints:	hierarchical execution and encapsulation	functions within a layer invoke (API of) others at next lower layer

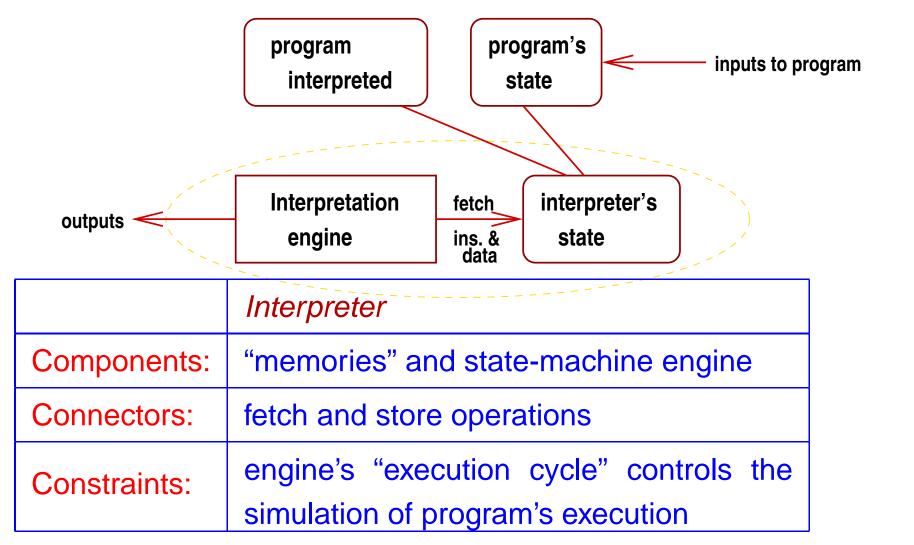
Examples: modular, object-oriented, N-tier systems (subroutine); communication protocols, operating systems (layered)



Advantages: hierarchical decomposition of solution; limits range of interactions between components, simplifying correctness reasoning; each layer defines a *virtual machine*; supports portability (by replacing lowest-level components).

Disadvantages: components must know the identities of other components to connect to them; side effects complicate correctness reasoning (e.g., A uses C, B uses and changes C, the result is an unexpected side effect from A's perspective; components sensitive to performance at lower levels/layers.

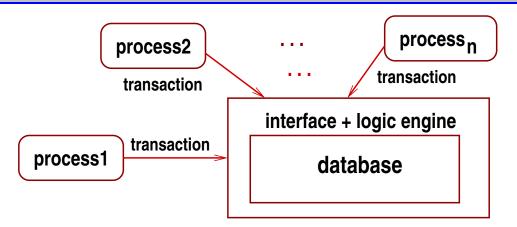
Virtual machine: interpreter



Examples: high-level programming-language interpreters, byte-code machines, virtual machines

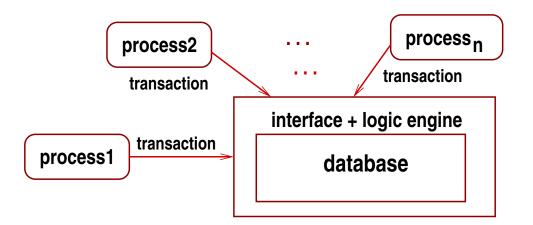
Advantages: rapid prototyping Disadvantages: inefficient.

Repositories: databases and blackboards



	Database	Blackboard
Components:	processes and database	knowledge sources and blackboard
Connectors:	queries and updates	notifications and updates
Constraints:	transactions (queries and updates) drive computation	knowledge sources respond when enabled by the state of the blackboard. Problem is solved by cooperative computation on blackboard.

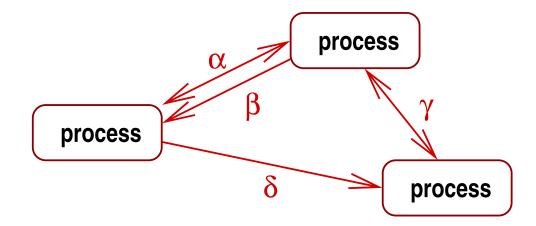
Examples: speech and pattern recognition (blackboard); syntax editors and compilers (parse tree and symbol table are repositories)



Advantages: easy to add new processes.

Disadvantages: alterations to repository affect all components.

Independent components: communicating processes

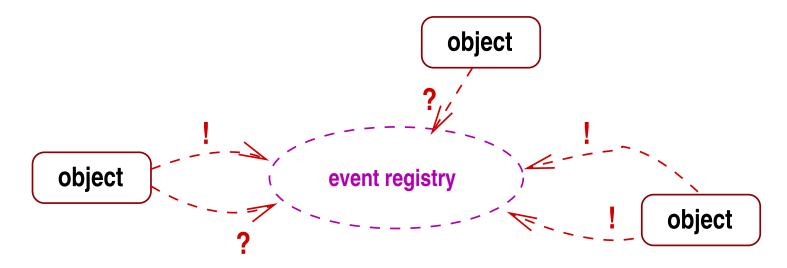


	Communicating processes
Components:	processes ("tasks")
Connectors:	ports or buffers or RPC
Constraints:	processes execute in parallel and send mes- sages (synchronously or asynchronously)

Example: client-server and peer-to-peer architectures

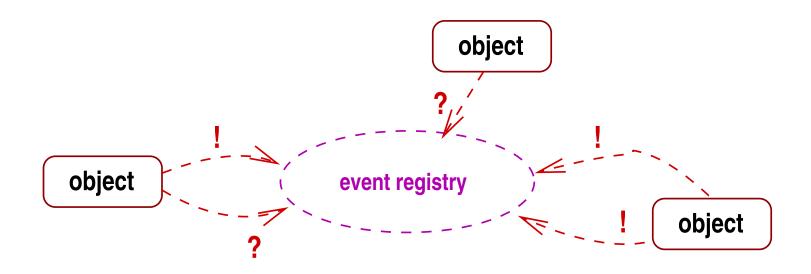
Advantages: easy to add and remove processes. *Disadvantages:* difficult to reason about control flow.

Independent components: event systems



	Event systems
Components:	objects or processes ("threads")
Connectors:	event broadcast and notification (implicit invocation)
Constraints:	components "register" to receive event notifi- cation; components signal events, environment notifies registered "listeners"

Examples: GUI-based systems, debuggers, syntax-directed editors, database consistency checkers



Advantages: easy for new listeners to register and unregister dynamically; component reuse.

Disadvantages: difficult to reason about control flow and to formulate system-wide invariants of correct behavior.

Three architectures for a compiler (Garlan and Shaw)

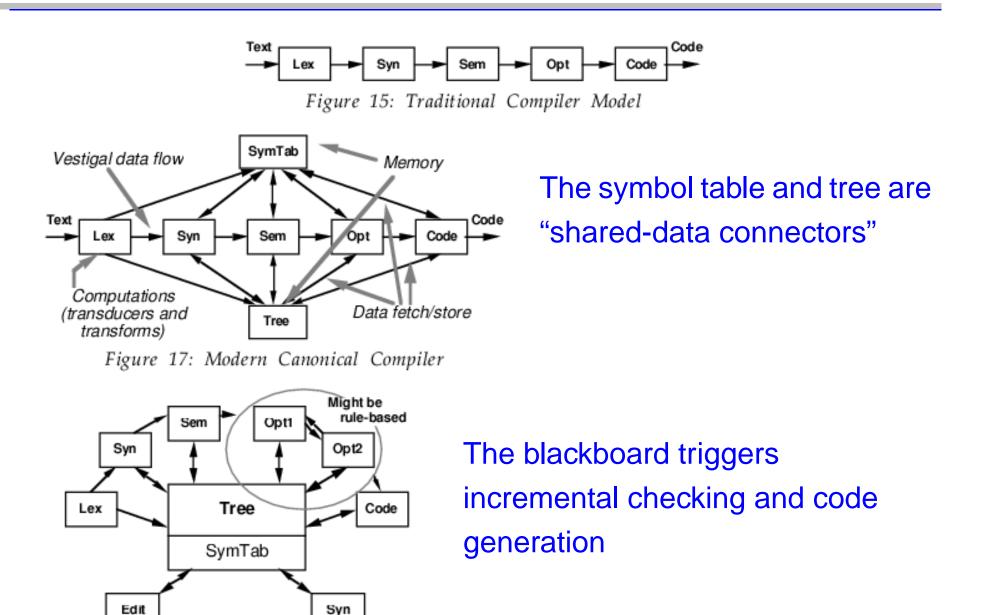


Figure 18: Canonical Compiler, Revisited

What do we gain from using a software architecture?

- the architecture helps us *communicate* the system's design to the project's stakeholders (users, managers, implementors)
- 2. it helps us *analyze* design decisions
- 3. it helps us *reuse* concepts and components in future systems

4. Architecture Description Languages

A language for connectors: UniCon

Shaw developed a language, *UniCon* (*Universal Connector Language*), for describing connectors and components.

Components are specified by **interfaces**, which include

(i) type;

(ii) attributes with values that specialize the type;

(iii) players, which are the component's connection points. Each player is itself typed.

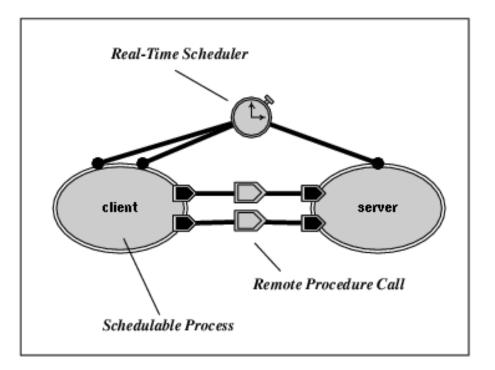
Connectors are specified by **protocols**; they have

(i) type;

(ii) specific properties that specialize the type;

(iii) roles that the connector uses to mediate (make) communication between components.

Graphical depiction of an assembly of three components and four connectors:



A development tool helps the designer draw the configuration and map it to coding.

Reference: M. Shaw, R. DeLine, and G. Zelesnik, Abstractions and Implementations for Architectural Connections. In 3d Int. Conf. Configurable Distributed Systems, Annapolis, Maryland, May 1996.

component Real_Time_System interface is type General end interface

implementation is uses client interface rtclient PRIORITY (10) ENTRYPOINT (client) end client

uses server interface rtserver PRIORITY (9) RPCTYPEDEF (new_type; struct; 12) RPCTYPESIN ("unicon.h") end server

establish RTM-realtime-sched with client.application1 as load client.application2 as load server.services as load ALGORITHM (rate_monotonic) PROCESSOR ("TESTBED.XX.EDU") TRACE (client.application1. external_interrupt1; client.application1.work_block1; server.services.work_block1; client.application1.work_block2; server.services.work_block2; client.application1.work_block3) TRACE (client.application2. external interrupt2; client.application2.work_block1; server.services.work_block1; client.application2.work_block2; server.services.work block2; client.application2.work_block3) end RTM-realtime-sched

establish RTM-remote-proc-call with client.timeget as caller server.timeget as definer IDLTYPE(Mach) end RTM-remote-proc-call

establish RTM-remote-proc-call with client-timeshow as caller server.timeshow as definer IDLTYPE(Mach) end RTM-remote-proc-call end implementation end Real. Time. System component RTClient interface is type SchedProcess PROCESSOR ("TESTBED.XX.EDU") TRIGGERDEF (external_interrupt1; 1.0) TRIGGERDEF (external_interrupt2; 0.5) SEGMENTDEF (work_block1; 0.02) SEGMENTDEF (work_block2; 0.03) SEGMENTDEF (work_block3; 0.05) player application1 is RTLoad TRIGGER (external_interrupt1) SEGMENTSET (work_block1, work_block2, work_block3) end application1 player application2 is RTLoad TRIGGER (external_interrupt2) SEGMENTSET (work_block1, work_block2, work_block3) end application2 player timeget is RPCCall SIGNATURE ("new_type *"; "void") end timeget player timeshow is RPCCall SIGNATURE ("void"; "void") end timeshow end interface

connector RTM-realtime-sched protocol is type RTScheduler role load is load end protocol

implementation is builtin end implementation end RTM-realtime-sched

connector RTM-remote-proc-call protocol is type RemoteProcCall role definer is definer role caller is caller end protocol

implementation is builtin end implementation end RTM-remote-proc-call *uses* statements instantiate the parts composed

connect statements state how players satisfy roles

bind statements map the external interface to the internal configuration

Connectors described in UniCon:

- data-flow connectors (pipe)
- procedural connectors (procedure call, remote procedure call): pass control
- data-sharing connectors (data access): export and import data
- resource-contention connectors (RT scheduler): competition for resources
- aggregate connectors (PL bundler): compound connections

Pipe Connector

Informal Description: The Unix abstraction for pipe, i.e. a bounded queue of bytes that are produced at a source and consumed at a sink. Also supports interactions between pipes and files, choosing the correct Unix implementation.

Icon: pipe section



Properties: PipeType, the kind of Unix pipe. Possible values Named, Unnamed

Roles: Source

Description: the source end of the pipe

Accepts player types: StreamOut of component Filter; ReadNext of component SeqFile Properties: MinConns, minimum number of connections. Integer values, default 1 MaxConns, maximum number of connections. Integer values, default 1

Sink

Description: the sink end of the pipe

Accepts player types: StreamIn of component Filter; WriteNext of component SeqFile Properties: MinConns, MaxConns, as for Source

ProcedureCall Connector

Informal Description: The architectural abstraction corresponding to the procedure call of standard programming languages. Requires signatures (eventually pre/post conditions) in the RoutineDef and RoutineCall players to match; if they don't, requests remediation. Supports renaming.

Icon: blunt arrowhead

Roles: Definer

 Description: role played by the procedure definition
 Accepts player types: RoutineDef of component Computation or Module
 Properties: MinConns, minimum number of definitions allowed. Integer, must be 1 MaxConns, maximum number of definitions allowed. Integer, must be 1

Caller

Description: the role played by the procedure call

Accepts player types: RoutineCall of component Computation or Module Properties: MinConns, minimum number of callers allowed. Integer, default 1 MaxConns, maximum number of callers allowed. Integer, default many

RemoteProcCall Connector

Informal Description: The abstraction for the remote procedure call facility supplied by the operating system. Requires signatures and eventually pre/post conditions in the RPCDef and RPCCall players to match. RemoteProcCall connectors require much more UniCon support than ProcedureCall connectors, as they must establish communication paths between processes.

Icon: bordered blunt arrowhead

Roles: Definer

Description: role played by the procedure definition Accepts player types: RPCDef of component Process or SchedProcess Properties: MinConns, MaxConns, as for ProcedureCall

Caller

Description: the role played by the procedure call Accepts player types: *RPCCall* of component Process or SchedProcess Properties: *MinConns, MaxConns*, as for ProcedureCall



DataAccess Connector

Informal Description: The architectural abstraction corresponding to imported and exported data of conventional programming languages.

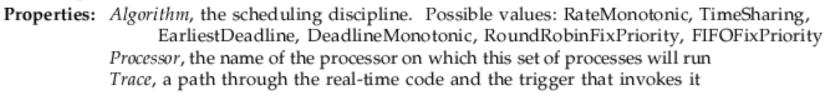
Icon: triangle

Roles: Definer, essentially similar to Definer of ProcedureCall User, essentially similar to Caller of ProcedureCall

RTScheduler Connector

Informal Description: Mediates competition for processor resources among a set of real-time processes (requires an operating system with appropriate real-time capabilities).

Icon: stopwatch



Roles: Load

Description: the role played by a real-time load on a processor Accepts player types: *RTLoad* of component SchedProcess

Properties: MinConns, minimum number of competing processes. Integer, default 2 MaxConns, maximum number of competing processes. Integer, default many



PLBundler Connector

Informal Description: A composite abstraction for matching definitions and uses of a collection of procedures and data. It allows multiple procedure and data definitions and uses to be matched with a single abstraction. Supports renaming.

Icon: chain links



Properties: Match, the correspondences between individual definitions in the bundles. Values are sets of pairs of names.

Roles: Participant

Description: a set of definitions and uses to take part in the linkage Accepts player types: *PLBundle* of component Computation, Module, or SharedData Properties: *MinConns*, minimum number of bundles to match. Integer, default 2 *MaxConns*, maximum number of bundles to match. Integer, default many Garlan and Allen developed Wright to specify protocols. Here is a single-client/single-server example:

System SimpleExample component Server = port provide [provide protocol] spec [Server specification] component Client = port request [request protocol] spec [Client specification] connector C-S-connector = role client [client protocol] role server [server protocol] glue [glue protocol]

Instances

s: Server

c: Client

cs: C-S-connector

Attachments

s.provide **as** cs.server; c.request **as** cs.client **end** SimpleExample.

The *protocols* are specified with Hoare's CSP (Communicating Sequential Processes) algebra.

connector C-S-connector =

role Client = (request!x → result?y → Client) □ § role Server = (invoke?x → return!y → Server) □ § glue = (Client.request?x → Server.invoke!x → Server.return?y → Client.result!y→glue) □ §

The *glue* protocol synchronizes the Client and Server roles:

 $\begin{array}{l} \texttt{Client} \parallel \texttt{Server} \parallel \texttt{glue} \\ \Rightarrow \texttt{result?y} \rightarrow \texttt{Client} \parallel \texttt{Server} \parallel \texttt{Server.invoke!x} \rightarrow \ldots \\ \Rightarrow \texttt{result?y} \rightarrow \texttt{Client} \parallel \texttt{return!y} \rightarrow \texttt{Server} \parallel \\ & \texttt{Server.return?y} \rightarrow \ldots \\ \Rightarrow \ldots \quad \Rightarrow \texttt{Client} \parallel \texttt{Server} \parallel \texttt{glue} \end{array}$

Forms of CSP processes:

- prefixing: $e \rightarrow P$
 - $plusOne?x \rightarrow return!x + 1 \rightarrow \cdots \parallel plusOne!2 \rightarrow return?y \rightarrow \cdots$
 - \Rightarrow return!2 + 1 $\rightarrow \cdots \parallel$ return?y $\rightarrow \cdots$
- external choice: P[Q

 $plusOne?x \rightarrow \cdots \text{ [] } plusTwo?x \rightarrow \cdots x + 2 \cdots \text{ || } plusTwo!5 \rightarrow \cdots$

- $\Rightarrow \cdots 5 + 2 \cdots \parallel \cdots$
- internal choice: $P \sqcap Q$

 $plusOne?x \rightarrow \cdots \parallel plusOne!5 \rightarrow \cdots \sqcap plusTwo!5 \rightarrow \cdots$

 \Rightarrow plusOne?x $\rightarrow \cdots \parallel$ plusTwo!5 $\rightarrow \cdots$

- parallel composition: P||Q
- halt: §
- (tail) recursion: p = · · · p (More formally, μz.P, where z may occur free in P.)

```
connector Pipe =
role Writer = write→Writer □ close→ §
role Reader = let ExitOnly = close→ §
in let DoRead = (read→Reader [] read-eof→ExitOnly)
in DoRead □ ExitOnly
glue = let ReadOnly = Reader.read→ReadOnly
[] Reader.read-eof →Reader.close → §
[] Reader.close→ §
in let WriteOnly = Writer.write→WriteOnly [] Writer.close→ §
in Writer.write→glue [] Reader.read→glue
[] Writer.close→ReadOnly [] Reader.close→WriteOnly
```

Reference: R. Allen and D. Garlan. A formal basis for architectural connection. *ACM TOSEM 1997.*

C2: an N-tier framework and language

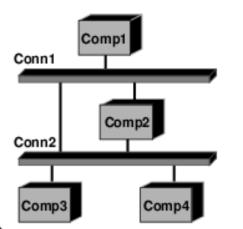
Developed at Univ. of California, Irvine, Institute of Software

Research: http://www.isr.uci.edu/architecture/c2.html

Examples of Domain- and Style-Specific Architectures

C2

- A component- and message-based style
 - for highly distributed software systems
- Generalized from GUI intensive systems' architectures
- C2 architectures are networks of concurrent components hooked together by connectors



- no component-to-component links
- "one up, one down" rule for components
- connector-to-connector links are allowed
- "many up, many down" rule for connectors
- all communication by exchanging messages
- substrate independence

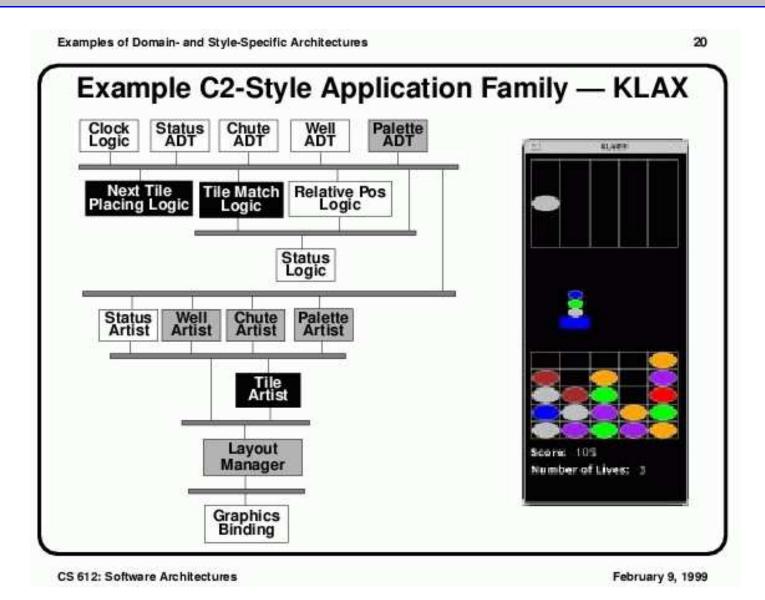
CS 612: Software Architectures

February 9, 1999

Diagrams are from Medvidovic's course,

http://sunset.usc.edu/classes/cs578_2002

Example architecture in C2: video game



Here is a **C2SADEL** description of the video game's "Well" component:

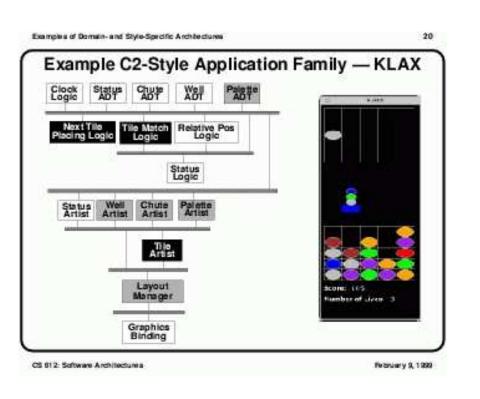
```
component WellADT is subtype Matrix (beh) {
   state {
      capacity : Integer;
      num tiles : Integer;
      well_at : Integer -> GSColor;
   invariant {
      (num_tiles \eqgreater 0) \and (num_tiles \eqless capacity);
   interface {
      prov gtl: GetTile (location : Integer) : Color;
      prov gt2: GetTile (i : Natural) : GSColor;
   operations {
      prov tileget: {
         let pos : Integer;
         pre (pos \greater 0) \and (pos \eqless num_tiles);
         post \result = well_at(pos) \and ~num_tiles = num_tiles - 1;
   }
   map
      gtl -> tileget (location -> pos);
      gt2 -> tileget (i -> pos);
}
```

Reference: N. Medvidovic, et al. A Language and Environment for Architecture-Based Software Development and Evolution. 21st Int. Conf. on Software Engineering, Los Angeles, May 1999.

And here is a description of a connector and part of the configuration:

```
message_filter no_filtering;
}
```

connector BroadcastConn is {



```
architectural_topology {
   component_instances {
     Well : WellADT;
     WellArt : WellArtist;
     MatchLogic : TileMatchLogic;
   }
   connector_instances {
     ADTConn : BroadcastConn;
     ArtConn : BroadcastConn;
   }
   connector ADTConn {
     top Well;
     bottom MatchLogic, ArtConn;
   }
   connector ArtConn {
     top ADTConn;
     bottom WellArt;
   }
}
```

ArchJava: Java extended with Unicon features

- Each component (class) has its own interfaces (*ports*) that list which methods it requires and provides
- Connectors are coded as classes, too, and extend the basic classes, Connector, Port, Method, etc.
- The ArchJava run-time platform includes a run-time type checker that enforces correctness of run-time connections (e.g., RPC, TCP)
- There is an open-source implementation and Eclipse plug-in
- www.archjava.org



```
package pos;
....
public component class POS {
    private final Sales sales = new Sales();
    private final UserInterface userInterface = new UserInterface();
    connect pattern Sales.model, UserInterface.view;
    connect pattern Sales.client, Inventory.server
       with TCPConnector {
       connect (Sales sender) throws Exception {
           return connect(sender.client, Inventory.server)
                    with new TCPConnector (connection, InetAddress.getByName(JDBC_SERVER_ADDRESS),
                                          JDBC_SERVER_PORT, JDBC_SERVER_NAME);
        ł
    1;
    public POS() {
       connect(sales.model, userInterface.view);
    }
    public void run() {
        sales.setData("Software Architecture in Practice, 2nd Edition");
    Ł
    public static void main (String[] args) {
        (new POS()).run();
    }
}
```

```
POS
                                       TCPconnector
               view
                          Sales C
  UserInterface
                                                       Inventory
                    model
                                 client
                                               server
package pos;
4.4.5
public component class Sales {
    private String data;
    public port model {
        provides String getData();
        provides void setData (String data);
        requires void updated();
    }
    public port interface client {
        requires connect() throws Exception;
        requires String executeUpdate (String statement);
    }
    public String getData() {
        return data;
    }
    public void setData(String data) {
        try {
            this.data = (new client()).executeUpdate(data);
        } catch (Exception e) {
            e. printStackTrace();
        model.updated();
    1
}
```

```
POS
                                  TCPconnector
             view
                       Sales
  UserInterface
                                                 Inventory
                  model
                             client
                                          server
package pos;
public component class Inventory {
    public port interface server {
      provides String executeUpdate (String statement);
    }
    public String executeUpdate (String statement) {
      return statement + " (validated)";
    }
    public Inventory () {
      try {
           TCPConnector.registerObject(this, POS.JDBC_SERVER_PORT,
                                         POS.JDBC_SERVER_NAME);
      } catch (IOException e) {
           e.printStackTrace();
       }
    }
    public static void main(String[] args) {
      new Inventory();
    }
}
```

From K. M. Hansen, www.daimi.dk/~marius/teaching/ATiSA2005



```
public class TCPConnector extends Connector {
```

```
// data members
protected TCPEndpoint endpoint;
// public interface
public TCPConnector(InetAddress host, int prt, String objName) throws IOException {
               endpoint = new TCPEndpoint(this, host, prt, objName);
ł
public Object invoke(Call call) throws Throwable {
               Method meth = call.getMethod();
               return endpoint.sendMethod(meth.getName(), meth.getParameterTypes(), call.getArguments());
public static void registerObject(Object o, int prt, String objName) throws IOException {
               TCPDaemon.createDaemon(prt).register(objName, o);
// interface used by TCPDaemon
TCPConnector(TCPEndpoint endpoint, Object receiver, String portName) {
               super(new Object[] { receiver }, new String[] { portName });
               this.endpoint = endpoint:
               endpoint.setConnector(this);
Object invokeLocalMethod(String name, Type parameterTypes[], Object arguments[]) throws Throwable {
               // find method with parameters that match parameterTypes
               Method meth = findMethod(name, parameterTypes);
               return meth.invoke(arguments);
}
```

So, what is an architectural description language?

It is a notation (linear or graphical) for specifying an architecture. It should specify

- structure: components (interfaces), connectors (protocols), configuration (both static and dynamic structure)
- behavior: semantical properties of individual components and connectors, patterns of acceptable communication, global invariants,
- design patterns: global constraints that support correctness-reasoning techniques, design- and run-time tool support, and implementation.

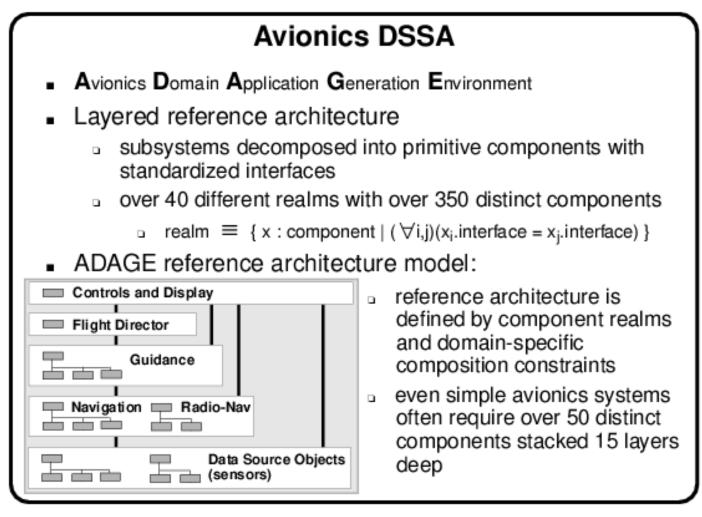
But it is difficult to design a *general-purpose* architectural description language that is *elegant*, *expressive*, and *useful*.

5. Domain-specific design

Domain-specific design

If the problem domain is a standard one (e.g., flight-control or telecommunications or banking), then there are precedents to follow.

- A Domain-Specific Software Architecture has
 - a domain: defines the problem area domain concepts and terminology; customer requirements; scenarios; configuration models (entity-relationship, data flow, etc.)
 - reference requirements: features that restrict solutions to fit the domain. ("Features" are studied shortly.) Also: platform, language, user interface, security, performance
 - ♦ a reference architecture
 - a supporting environment/infrastructure: tools for modelling, design, implementation, evaluation; run-time platform
 - a process or methodology to implement the reference architecture and evaluate it.



CS 612: Software Architectures

February 2, 1999

from Medvidovic's course, http://sunset.usc.edu/classes/cs578_2002

Domain-specific (modelling) language (DSL)

is a modelling language specialized to a specific problem domain, e.g., telecommunications, banking, transportation.

We use a DSL to describe a problem and its solution in concepts familiar to people who work in the domain.

It might define (entity-relationship) models, ontologies (class hierarchies), scenarios, architectures, and implementations.

Example: a DSL for sensor-alarm networks: *domains:* sites (building, floor, hallway, room), devices (alarm, movement detector, camera, badge), people (employee, guard, police, intruder). Domain elements have *features/attributes* and *operations. Actions* can be by initiated by *events* — "when a movement detector detects an intruder in a room, it generates a movement-event for a camera and sends a message to a guard...."

When a DSL can generate computer implementations, it is a *domain-specific programming language*.

Domain-specific programming language

In the Unix world, these are "little languages" or "mini-languages," designed to solve a specific class of problems. Examples are awk, make, lex, yacc, ps, and Glade (for GUI-building in X).

Other examples are Excel, HTML, XML, SQL, regular-expression notation and BNF. These are called *top-down* DSLs, because they are designed to implement domain concepts and nothing more. Non-programmers can use a top-down DSL to write solutions.

The *bottom-up* approach, called *embedded* or *in-language DSL*, starts with a dynamic-data-structure language, like Scheme or Perl or Python, and adds libraries (modules) of functions that encode domain-concepts-as-code, thus "building the language upwards towards the problem to be solved." Experienced programmers use bottom-up DSLs to program solutions.

Tradeoffs in using (top-down) DSLs

- non-programmers can discuss and use the DSL
- the DSL supports patterns of design, implementation, and optimization
- ✓ fast development
- **X** staff must be trained to use the DSL
- interaction of DSL-generated software with other software components can be difficult
- **×** there is high cost in developing and maintaining a DSL

Reference: J. Lawall and T. Mogensen. Course on Scripting Languages and DSLs, Univ. Copenhagen, 2006, www.diku.dk/undervisning/2006f/213

6. Software product lines

A software product line

is also called a *software system family* — a collection of software products that share an architecture and components, constructed by a product line. They are inspired by the products produced by industrial assembly lines, e.g., automobiles.

The CMU Software Engineering Institute definition:

A product line is a set of software intensive systems that (i) share a common set of features, (ii) satisfy the needs of a particular mission, and (iii) are developed from a set of core assets in a prescribed way.

Key issues:

variability: Can we state precisely the products' variations (*features*) ? *guidance:* Is there a precise recipe that guides feature selection and product assembly?

Reference: www.softwareproductlines.com

An example product line: Cummins Corporation

produces diesel engines for trucks and heavy machinery. An engine controller has 100K-200K lines-of-code. At level of 12 engine "builds," company switched to a product line approach:

- 1. defined engine controller domain
- 2. defined a reference architecture
- 3. built reusable components
- 4. required all teams to follow product line approach

Cummins now produces 20 basic "builds" — 1000 products total; development time dropped from 250 person/months to < 10. A new controller consists of 75% reused software.

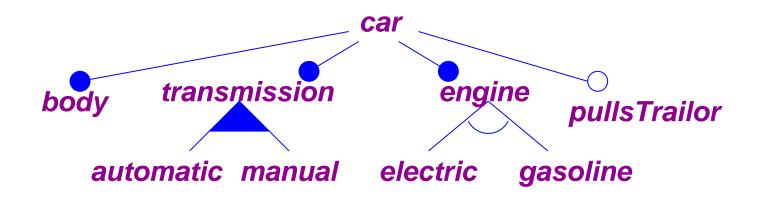
Reference: S. Cohen. Product line practice state of the art report. CMU/SEI-2002-TN-017. are a development tool for domain-specific architectures and product lines. They help define a domain's reference requirements and guide implementions of instances of the reference architecture.

A *feature* is merely a property of the domain. (Example: the features/options/choices of an automobile that you order from the factory.)

A *feature diagram* displays the features and guides a user in choosing features for the solution to a domain problem.

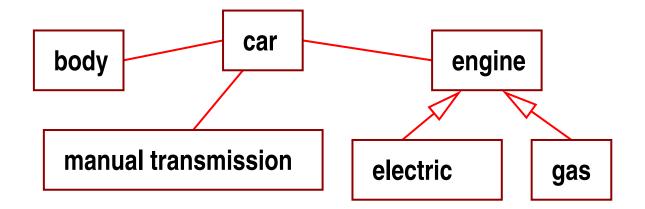
It is a form of decision tree with *and-or-xor* branching, and its hierarchy reflects dependencies of features as well as modification costs.

Feature diagram for assembling automobiles

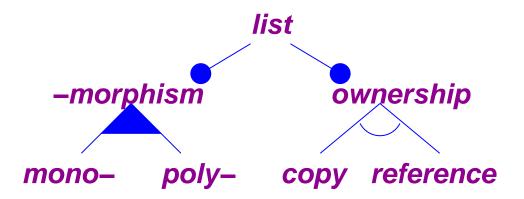


Filled circles label required features; unfilled circles label optional ones. Filled arcs label xor-choices; unfilled arcs label or-choices (where at least one choice is selected).

Here is one possible outcome of "executing" the feature diagram:



Feature diagrams work well for configuring generic data structures:



Compare the diagram to the typical class-library representation of a generic list structure.

An advantage of a feature-diagram construction of a list structure over a class-library construction is that the former can generate a smaller, more efficient list structure, customized to exactly the choices of the client. Feature diagrams are useful for both *constraining* as well as *generating* an architecture: the feature requirements are displayed in a feature diagram, which guides the user to generating the desired instance of the reference architecture.

Feature diagrams are an attempt at making software assembly appear similar to assembly of mass-produced products like automobiles.

In particular, feature diagrams encourage the use of *standardized, parameterized, reusable software components*.

Feature diagrams might be implemented by a tool that selects components according to feature selection. Or, they might be implemented within the structure of a *domain-specific programming language* whose programs select and assemble features.

Reference: K. Czarnecki and U. Eisenecker. *Generative Programming.* Addison-Wesley 2000. is the name given to the application of programs that generate other programs (cf. "automatic programming" in the 1950s). A compiler is of course a generating program, but so are feature-diagram-driven frameworks, partial evaluators, and some development environments (e.g., for Java beans).

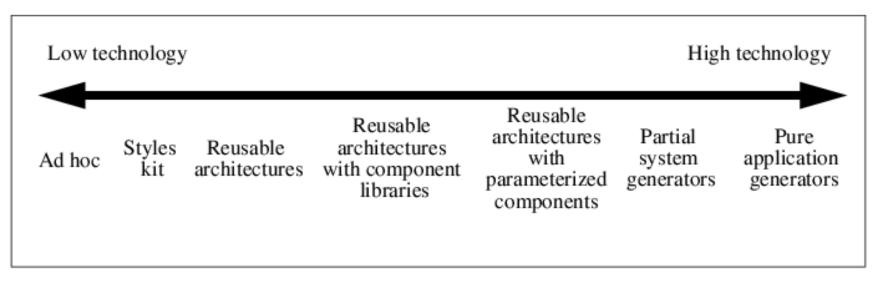


Figure 1: Technology spectrum for architecture selection and creation

Reference: Coming attractions in software architecture, P. Clements. CMU/SEI-96-TR-008.

Generative programming is motivated by the belief that conventional software production methods (even those based on "object-oriented" methodologies) will never support component reuse:

graphical user interface					
web browser				etc.	
browser					
core competence of software development organization					
		1	Corba/	i	
en ayption/	etc.		Com/	web	
decryption			Beans	server	
database management system					
operating system					

software product

Reference: Jan Bosch. Design and Use of Software Architectures. Addison-Wesley, 2000.

One solution is to understand a software system as a customized product, produced by generative programming, from a product line. Reference: K. Czarnecki and U. Eisenecker. *Generative Programming.* Addison-Wesley 2000.

10. Final Remarks

TABLE 1. Academic versus industrial view on software architecture

Academia

- Architecture is explicitly defined.
- Architecture consists of components and first-class connectors.
- Architectural description languages (ADLs) explicitly describe architectures and are used to automatically generate applications.

Industry

- Mostly conceptual understanding of architecture. Minimal explicit definition, often through notations.
- No explicit first-class connectors (sometimes ad-hoc solutions for run-time binding and glue code for adaptation between assets).
- Programming languages (e.g., C++) and script languages (e.g., Make) used to describe the configuration of the complete system.

Reference: Jan Bosch. *Design and Use of Software Architectures*. Addison-Wesley, 2000.

Academia	Industry		
 Reusable components are black-box entities. 	 Components are large pieces of software (sometimes more than 80 KLOC) with a complex internal structure and no enforced encapsulation boundary, e.g., object-oriented frameworks. 		
 Components have nar- row interface through a single point of access. 	 The component interface is provided through entities, e.g., classes in the component. These interface entities have no explicit differences to non-interface entities. 		
 Components have few and explicitly defined variation points that are configured during instantiation. 	 Variation is implemented through configura- tion and specialization or replacement of entities in the component. Sometimes multi- ple implementations (versions) of compo- nents exist to cover variation requirements 		
 Components imple- ment standardized interfaces and can be traded on component markets. 	 Components are primarily developed inter- nally. Externally developed components go through considerable (source code) adapta- tion to match the product-line architecture requirements. 		
 Focus is on component functionality and on the formal verification of functionality. 	 Functionality and quality attributes, e.g. per- formance, reliability, code size, reusability and maintainability, have equal importance. 		

TABLE 2. Academic versus industrial view on reusable components

Selected textbook references

F. Buschmann, et al. *Pattern-Oriented Software Architecture.* Wiley 1996.

P. Clements and L. Northrup. *Software Product Lines*. Addison-Wesley 2002.

P. Clements, et al. *Documenting Software Architectures: Views and Beyond*. Addison Wesley, 2002.

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E. Gamma, et al. *Design Patterns: Elements of Reusable Object-Oriented Software*. Addison Wesley, 1994.

M. Shaw and D. Garlan. Software Architecture. Prentice Hall 1996.